

Projection

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Model

$$\begin{aligned}c_t^{-\nu} &= \mathbf{E}_t \left[\beta c_{t+1}^{-\nu} \alpha z_{t+1} k_{t+1}^{\alpha-1} \right] \\c_t + k_{t+1} &= z_t k_t^\alpha \\ \ln(z_{t+1}) &= \rho \ln(z_t) + \varepsilon_{t+1} \\ &\varepsilon_{t+1} \sim N(0, \sigma^2) \\ &k_1, z_1 \text{ given}\end{aligned}$$

Projection Methods

True rational expectations solution:

$$\begin{aligned}c_t &= c(k_t, z_t) \\ k_{t+1} &= k(k_t, z_t)\end{aligned}$$

- Why a difficult problem to find these?

Define error terms

$$e(k_t, z_t) = -c_t^{-\nu} + \mathbb{E}_t \left[\beta c_{t+1}^{-\nu} \alpha z_{t+1} k_{t+1}^{\alpha-1} \right]$$

At the true solutions, $c(k_t, z_t)$ and $k(k_t, z_t)$:

$$e(k_t, z_t) = 0 \quad \forall k_t, z_t$$

- Structural parameters $(\alpha, \beta, \rho, \sigma)$ have fixed numerical values (thus not included as arguments in policy function)

$$c_t = c(k_t, z_t) \approx P_n(k_t, z_t; \eta_n)$$

- $P_n(\cdot)$: from class of approximating functions
 - such as polynomials or splines
 - n is fixed \implies solve for η_n , a *finite-dimensional* object

Which equations to use?

- goal: solve for $P_n(k_t, z_t; \eta_n) \approx c(k_t, z_t)$,
 - i.e., N_n elements of η_n
 - $k(k_t, z_t)$ implicitly defined by budget constraint
- One first-order equation left, namely Euler equation
 - this is a different equation at each point in the state space
 - \implies plenty of equations

Which equations to use?

- At M grid points $\{k_i, z_i\}$ with $M \geq N_n$ we would like the following to equal zero:

$$e(k_i, z_i; \eta_n) = -P_n(k_i, z_i; \eta_n)^{-\nu} +$$

$$E \begin{bmatrix} \alpha\beta \times \\ P_n(\{\mathbf{k}'\}, \{\mathbf{z}'\}; \eta_n)^{-\nu} \times \\ \{\mathbf{z}'\} \times \\ (\{\mathbf{k}'\})^{\alpha-1} \end{bmatrix}$$

Which equations to use?

- **Goal:** \forall grid point get an expression with η_n as only unknown

$$e(k_i, z_i; \eta_n) = -P_n(k_i, z_i; \eta_n)^{-\nu} +$$

$$E \begin{bmatrix} \alpha\beta \times \\ P_n(\mathbf{k}', \mathbf{z}'; \eta_n)^{-\nu} \times \\ \mathbf{z}' \times \\ (\mathbf{k}')^{\alpha-1} \end{bmatrix}$$

- Note that k_i and z_i are known

Which equations to use?

$$\begin{aligned}
 e(k_i, z_i; \eta_n) = & -P_n(k_i, z_i; \eta_n)^{-\nu} + \\
 & \alpha\beta \times \\
 \text{E} \left[& P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n), \exp\{\rho \ln(z_i) + \varepsilon'\}; \eta_n)^{-\nu} \times \right. \\
 & \exp\{\rho \ln(z_i) + \varepsilon'\} \times \\
 & \left. (z_i k_i^\alpha - P_n(k_i, z_i; \eta_n))^{\alpha-1} \right]
 \end{aligned}$$

How to deal with expectations operator?

Let $\{\omega_j, \zeta_j\}_{j=1}^J$ be the Hermite Gaussian quadrature nodes

$$e(k_i, z_i; \eta_n) = -P_n(k_i, z_i; \eta_n)^{-\nu} +$$

$$\alpha\beta \times$$

$$P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n), \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n)^{-\nu} \times$$

$$\sum_{j=1}^J$$

$$\exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times$$

$$(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n))^{\alpha-1}$$

$$\omega_j / \sqrt{\pi}$$

Define error terms

$$e(k_i, z_i; \eta_n) = -P_n(k_i, z_i; \eta_n)^{-\nu} +$$

$$\alpha\beta \times$$

$$P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n), \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n)^{-\nu} \times$$

$$\exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times$$

$$(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n))^{\alpha-1}$$

$$\omega_j / \sqrt{\pi}$$

$$\sum_{j=1}^J$$

How to find coefficients of approximation?

- True rational expect. solution gives zero error term $\forall(k_i, z_i)$
- Thus, choose η_n such that error terms are as small as possible.
- **Collocation** ($M = N_n$): Use equation solver to get errors exactly equal to zero on grid
- **Galerkin** ($M > N_n$): Use minimization routine (and possibly smart weighting of error terms)

Different types of approximating functions

- $P_n(k_i, z_i; \eta_n)$ could be polynomial or spline
- dimension η_n usually higher for splines
 - may make eq. solver/minimization less appropriate
 - use iteration scheme instead

How to find coefficients of approximation?

- 1 Equation solver or minimization routine
- 2 Iteration procedures
 - 1 fixed-point iteration
 - 2 time iteration

Iterating versus eq. solver/minimization

- Advantage:
 - less of a black box
 - can deal with many coefficients
 - e.g. when spline is used
 - some iteration schemes are guaranteed to converge
 - under some regularity conditions
- Disadvantage:
 - does not use information on how best to update

Iteration procedure: Construct Grid

- Construct a grid with nodes for k and z
- At the nodes construct the basis functions of $P_n(k, z; \eta_n)$.
- For example, if

$$P_n(k, z; \eta_n) = \eta_{0,n} + \eta_{k,n}k + \eta_{z,n}z + \eta_{kk}k^2 + \eta_{kz}kz + \eta_{zz}z^2$$

then construct the matrix (where subscripts denote grid numbers)

$$X = \begin{bmatrix} 1 & k_1 & z_1 & k_1^2 & k_1 z_1 & z_1^2 \\ 1 & k_2 & z_2 & k_2^2 & k_2 z_2 & z_2^2 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & k_M & z_M & k_M^2 & k_M z_M & z_M^2 \end{bmatrix}$$

and calculate $(X'X)^{-1} X'$

Iteration procedure: Construct Grid

- **Chebyshev nodes:** Using Chebyshev nodes is important. This ensures uniform convergence. With equidistant nodes it is possible that the oscillations between grid point explode as the order of the polynomial increases.
- **Chebyshev polynomials:** If you have (i) no problems finding initial conditions and (ii) only low-order approximations so that calculating the inverse of $X'X$ can be done accurately, then you can use regular polynomials. Orthogonal Chebyshev polynomials can overcome these problems. They ensure that $X'X$ is diagonal (and trivial to invert). This does require scaling of the state variables so they are between -1 and 1 .

Fixed-point Iteration

The value of η_n used in the q^{th} iteration is referred to as η_n^q . Follow the following iteration scheme until convergence

- At each grid point:
 - Calculate the RHS of the Euler equation using the latest value for η_n , i.e., η_n^{q-1}
 - Use RHS to calculate c_i , value for c at i^{th} grid point
- Use values for c_i to obtain an estimate for η_n , $\hat{\eta}_n^q$
 - Polynomial: run a regression to get $\hat{\eta}_n^q$
 - Spline: the values of c at the nodes are the new values of η_n
- Let $\eta_n^q = \lambda \hat{\eta}_n^q + (1 - \lambda) \eta_n^{q-1}$

Fixed-point Iteration

- **Step 1: Calculate current consumption values implied by η_n^{j-1} at each grid point**
 - Use η_n^{q-1} to calculate $k' = z_i k_i^\alpha - P_n(k_i, z_i; \eta_n^{q-1})$
 - Use η_n^{q-1} to calculate $c' = P_n(k', z'; \eta_n^{q-1})$
 - Then, get c_i from

$$(c_i)^{-\nu} =$$

$$\sum_{j=1}^J \left[\begin{array}{c} \alpha\beta \times \\ P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n^{q-1}), \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n^{q-1})^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ \left(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n^{q-1}) \right)^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

Fixed-point iteration

Step 2: Get new estimate for η_n by running a projection step

- Let $Y = [c_1, c_2, \dots, c_M]'$
- If

$$P_n(k, z; \eta_n) = \eta_{0,n} + \eta_{k,n}k + \eta_{z,n}z + \eta_{kk}k^2 + \eta_{kz}kz + \eta_{zz}z^2$$

then

$$\hat{\eta}_n^q = (X'X)^{-1} X'Y$$

Fixed-point iteration

Step 2: Get new estimate for η_n by running a projection step

- If

$$P_n(k, z; \eta_n) = \exp \left(\eta_{0,n} + \eta_{k,n}k + \eta_{z,n}z + \eta_{kk}k^2 + \eta_{kz}kz + \eta_{zz}z^2 \right)$$

then

$$\hat{\eta}_n^q = (X'X)^{-1} X' \ln(Y)$$

- no stochastic error term \implies ok to take \ln of LHS & RHS

Fixed-point iteration

Step 3: Update η_n

$$\eta_n^q = \lambda \widehat{\eta}_n^q + (1 - \lambda) \eta_n^{q-1} \quad \text{for } 0 < \lambda \leq 1$$

- Fixed-point iteration does not always converge
 - Choosing a lower value of λ :
 - convergence more likely
 - slows down algorithm if lower value not needed for convergence
- Alternative is **time iteration**

Time Iteration

- At each grid point use η_n^{q-1} only for *next period's* choices
- Again solve for c_i at each grid point
 - this is now a bit trickier (non-linear problem)
- Get n_n^q as with fixed-point iteration
 - guaranteed to converge without dampening (under regularity conditions)

Time Iteration - solving for c

Solve c_i from following non-linear equation

$$(c_i)^{-\nu} = \sum_{j=1}^J \left[\begin{array}{c} \alpha\beta \times \\ P_n(z_i k_i^\alpha - c_i, \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n^{q-1})^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ (z_i k_i^\alpha - c_i)^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

Time Iteration

- Natural interpretation for η_n^{q-1} and η_n^q , namely
 - η_n^{q-1} is tomorrow's policy function and
 - η_n^q is today's policy function
- Time iteration is reliable and convergent
 - (the proof is related to the convergence of value function iteration, which uses the same idea)

Fixed-point versus time iteration

- Fixed-point iteration uses η_n^{q-1} for *all* terms on the RHS, i.e., both next period's consumption choice and today's capital choice
- Time iteration uses η_n^{q-1} only to evaluate next period's consumption
- The structure of time iteration mimics the choice of value function iteration:
 - next period's behavior described by previous solution for value function
 - Bellman equation used to solve for choice of c and k *simultaneously*

Endogenous grid points

- Simple idea: construct grid for k' instead of a grid for k
- Instead of solving for the choice k' given k , we now solve for the value of k that would have led to the choice k'
- In both cases you end up at each grid point with a set of values for k and a set of corresponding values for k' .
- Terminology is a bit confusing: the grid itself is exogenous and fixed but it is for an endogenous variable
- You can use endogenous grid points both with fixed-point and with time iteration
- The added value with time iteration lies in getting rid of the non-linear problem of solving for today's choices

Endogenous grid points and time iteration

- Time iteration \implies
 - use η_n^{q-1} for tomorrow's choices and
 - use η_n^q only for today's choices (which show up on both sides of the policy function)
- Then, get c_i from

Endogenous grid points and time iteration

$$(c_i)^{-\nu} = \sum_{j=1}^J \left[\begin{array}{c} \alpha\beta \times \\ P_n(k'_i, \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n^{q-1})^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ (k'_i)^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

and k_i from

$$k'_i + c_i = z_i k^\alpha$$

Perturbation versus projection

- Nondifferentiabilities
 - impossible for perturbation
- Large number of state variables
 - difficult for projection
- Constructing the grid can be difficult
 - apriori hard to know what sensible points are
 - some calculations may not be well defined everywhere

Perturbation versus projection

- Global versus local
 - Projection designed to be global method
 - Perturbation designed to be local method
 - but could give accurate global approximation
 - question is whether (lower-order) derivatives at perturbation point capture global behavior

When can't you use projection methods?

- Not all solutions to optimization problems can be characterized by first-order conditions
 - e.g. when objective function is not concave or budget set not convex
 - then you have no choice but to use Value Function Iteration

When can't you use projection methods?

- Constructing a grid where all calculations are well defined may be tough
 - e.g., not get negative consumption/unemployment
 - this can be tough even at the true solution
 - calculations should be possible also on path towards solution
- Solutions
 - Simply exclude problematic grid points (works for Galerkin)
 - Endogenize grid using simulations (Parameterized expectations)
 - but simulated points cluster so you are likely to get worse convergence properties

References

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 - shows that time-iteration converges even in the presence of inequality constraints